

## Codes over $F_{p^2}$ and $F_p \times F_p$ , lattices, and theta functions

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Let  $\ell > 0$  be a square free integer and  $\mathcal{O}_K$  the ring of integers of the imaginary quadratic field  $K = Q(\sqrt{-\ell})$ . Codes  $C$  over  $K$  determine lattices  $\Lambda_\ell(C)$  over rings  $\mathcal{O}_K/p\mathcal{O}_K$ . The theta functions  $\theta_{\Lambda_\ell}(C)$  of such lattices are known to determine the symmetrized weight enumerator  $swe(C)$  for small primes  $p = 2, 3$ ; see [1], [10].

In this paper we explore such constructions for any  $p$ . If  $p \nmid \ell$  then the ring  $\mathcal{R} := \mathcal{O}_K/p\mathcal{O}_K$  is isomorphic to  $\mathbb{F}_{p^2}$  or  $\mathbb{F}_p \times \mathbb{F}_p$ . Given a code  $C$  over  $\mathcal{R}$  we define new theta functions on the corresponding lattices. We prove that the theta series  $\theta_{\Lambda_\ell}(C)$  can be written in terms of the complete weight enumerator of  $C$  and that  $\theta_{\Lambda_\ell}(C)$  is the same for almost all  $\ell$ . Furthermore, for large enough  $\ell$ , there is a unique complete weight enumerator polynomial which corresponds to  $\theta_{\Lambda_\ell}(C)$ .

*Keywords:* codes, lattices, theta functions

### 1. Introduction

Let  $\ell > 0$  be a square free integer,  $K = Q(\sqrt{-\ell})$  be the imaginary quadratic field, and  $\mathcal{O}_K$  its ring of integers. Codes, Hermitian lattices, and their theta-functions over rings  $\mathcal{R} := \mathcal{O}_K/p\mathcal{O}_K$ , for small primes  $p$ , have been studied by many authors, see [7, 8], [1] among others. In [1], an explicit description of theta functions and MacWilliams identities are given for  $p = 2, 3$ . For a general reference of the topic, see [6].

In this paper we aim to explore such constructions, under certain restric-

tions, for any  $p$ . Further, we study the weight enumerators of such codes in terms of the theta functions of the corresponding lattices. We aim to find MacWilliams-like identities in such cases and explore to what extent the theta functions of these lattices determine the codes. The last question was studied in [2] and [10] for  $p = 2$ .

This paper is organized as follows. In section 2 we give a brief overview of the basic definitions for codes and lattices and define theta functions over  $\mathbb{F}_p$ . In section 3 we define theta-functions on the lattice defined over  $\mathcal{R} := \mathcal{O}_K/p\mathcal{O}_K$ . For general odd  $p$ , among the  $p^2$  lattices, there are  $\frac{(p+1)^2}{4}$  associated theta series.

In section 4, we address a special case of a general problem of the construction of lattices: the injectivity of Construction A. For codes defined over an alphabet of size four (regarded as a quotient of the ring of integers of an imaginary quadratic field), the problem is solved completely in [10]. The analogous questions are asked for codes defined over  $\mathbb{F}_{p^2}$  or  $\mathbb{F}_p \times \mathbb{F}_p$ . The main obstacle seems to express the theta function in terms of the symmetric weight enumerator of the code. However, the theta function  $\theta_{\Lambda_\ell}(C)$  can be expressed in terms of the complete weight enumerator of the code (cf. section 4). Using such an expression we prove the following two facts:

**Theorem:** Let  $p$  be a fixed prime and  $\ell$  any square free integer such that  $K = \mathbb{Q}(\sqrt{-\ell})$  and  $\mathcal{R} := \mathcal{O}_K/p\mathcal{O}_K$  is isomorphic to  $\mathbb{F}_{p^2}$  or  $\mathbb{F}_p \times \mathbb{F}_p$ . For a given code  $C$  defined over  $\mathcal{R}$ , the theta series  $\theta_{\Lambda_\ell}(C)$  is the same for almost all  $\ell$ .

**Theorem:** Let  $C$  be a code of size  $n$  defined over  $\mathcal{R}$  and  $\theta_{\Lambda_\ell}(C)$  be its corresponding theta function for level  $\ell$ . Then, for large enough  $\ell$ , there is a unique complete weight enumerator polynomial which corresponds to  $\theta_{\Lambda_\ell}(C)$ .

In contrary to results in [10] we did not attempt to find explicit bounds for  $\ell$ . However, for a given small  $p$  it is possible such bounds can be determined using similar techniques as in [10]. This is intended to be completed in further work; see [11].

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## 2. Preliminaries

Let  $\ell > 0$  be a square free integer and  $K = Q(\sqrt{-\ell})$  be the imaginary quadratic field with discriminant  $d_K$ . Recall that

$$d_K = \begin{cases} -\ell & \text{if } \ell \equiv 3 \pmod{4}, \\ -4\ell & \text{otherwise.} \end{cases}$$

Let  $\mathcal{O}_K$  be the ring of integers of  $K$ . A lattice  $\Lambda$  over  $K$  is an  $\mathcal{O}_K$ -submodule of  $K^n$  of full rank. The Hermitian dual is defined by

$$\Lambda^* = \{x \in K^n \mid x \cdot \bar{y} \in \mathcal{O}_K, \text{ for all } y \in \Lambda\}, \quad (1)$$

where  $x \cdot y := \sum_{i=1}^n x_i y_i$  and  $\bar{y}$  denotes component-wise complex conjugation. In the case that  $\Lambda$  is a free  $\mathcal{O}_K$ -module, for every  $\mathcal{O}_K$  basis  $\{v_1, v_2, \dots, v_n\}$  we can associate a Gram matrix  $G(\Lambda)$  given by  $G(\Lambda) = (v_i \cdot v_j)_{i,j=1}^n$  and the determinant  $\det \Lambda := \det(G)$  defined up to squares of units in  $\mathcal{O}_K$ . If  $\Lambda = \Lambda^*$  then  $\Lambda$  is Hermitian self-dual (or unimodular) and integral if and only if  $\Lambda \subset \Lambda^*$ . An integral lattice has the property  $\Lambda \subset \Lambda^* \subset \frac{1}{\det \Lambda} \Lambda$ . An integral lattice is called even if  $x \cdot x \equiv 0 \pmod{2}$  for all  $x \in \Lambda$ , and otherwise it is odd. An odd unimodular lattice is called a Type 1 lattice and even unimodular lattice is called a Type 2 lattice.

The theta series of a lattice  $\Lambda$  in  $K^n$  is given by

$$\theta_\Lambda(\tau) = \sum_{z \in \Lambda} e^{\pi i \tau z \cdot \bar{z}},$$

where  $\tau \in H = \{z \in \mathbb{C} : \text{Im}(z) > 0\}$ . Usually we let  $q = e^{\pi i \tau}$ . Then,  $\theta_\Lambda(q) = \sum_{z \in \Lambda} q^{z \cdot \bar{z}}$ . The one dimensional theta series (or Jacobi's theta series) and its shadow are given by

$$\theta_3(q) = \sum_{n \in \mathbb{Z}} q^{n^2}, \quad \theta_2(q) = \sum_{n \in \frac{1}{2} + \mathbb{Z}} q^{n^2}.$$

Let  $\ell \equiv 3 \pmod{4}$  and  $d$  be a positive number such that  $\ell = 4d - 1$ . Then,  $-\ell \equiv 1 \pmod{4}$ . This implies that the ring of integers is  $\mathcal{O}_K = \mathbb{Z}[\omega_\ell]$ , where  $\omega_\ell = \frac{-1 + \sqrt{-\ell}}{2}$  and  $\omega_\ell^2 + \omega_\ell + d = 0$ . The principal norm form of  $K$  is given by

$$Q_d(x, y) = |x - y\omega_\ell|^2 = x^2 + xy + dy^2. \quad (2)$$

The structure of  $\mathcal{O}_K/p\mathcal{O}_K$  depends on the value of  $\ell$  modulo  $p$ . For  $\left(\frac{a}{p}\right)$  the Legendre symbol,

$$\mathcal{O}_K/p\mathcal{O}_K = \begin{cases} \mathbb{F}_p \times \mathbb{F}_p & \text{if } \left(\frac{-\ell}{p}\right) = 1, \\ \mathbb{F}_{p^2} & \text{if } \left(\frac{-\ell}{p}\right) = -1, \\ \mathbb{F}_p + u\mathbb{F}_p \text{ with } u^2 = 0 & \text{if } p \mid \ell. \end{cases} \quad (3)$$

We will concern ourselves with the cases where  $p \nmid \ell$ .

### 2.1. Theta functions over $\mathbb{F}_p$

Let  $q = e^{\pi i \tau}$ . For integers  $a$  and  $b$  and a prime  $p$ , let  $\Lambda_{a,b}$  denote the lattice  $a - b\omega_\ell + p\mathcal{O}_K$ . The theta series associated to this lattice is

$$\begin{aligned} \theta_{\Lambda_{a,b}}(q) &= \sum_{m,n \in \mathbb{Z}} q^{|a+mp-(b+np)\omega_\ell|^2} \\ &= \sum_{m,n \in \mathbb{Z}} q^{Q_d(mp+a, np+b)} \\ &= \sum_{m,n \in \mathbb{Z}} q^{p^2 Q_d(m+a/p, n+b/p)}. \end{aligned} \quad (4)$$

For a prime  $p$  and an integer  $j$ , consider the theta series

$$\theta_{p,j}(q) := \sum_{n \in \frac{j}{2p} + \mathbb{Z}} q^{n^2}. \quad (5)$$

Note that  $\theta_{p,j}(q) = \theta_{p,k}(q)$  if and only if  $j \equiv \pm k \pmod{2p}$ .

The theta series of  $\Lambda_{a,b}$  can be written in terms of these series. In particular,

$$\theta_{\Lambda_{a,b}}(q) = \theta_{p,b}(q^{p^2})\theta_{p,2a+b}(q^{p^2}) + \theta_{p,b+p}(q^{p^2})\theta_{p,2a+b+p}(q^{p^2}). \quad (6)$$

The proof of this fact is similar to the proof of Lemma 2.1 in [5].

**Lemma 2.1.** *For any integers  $a, b, m, n$ , if the ordered pair  $(m, n)$  is component-wise congruent modulo  $p$  to one of*

$$(a, b), (-a, -b), (a + b, -b), (-a - b, b),$$

then

$$\theta_{\Lambda_{m,n}}(q) = \theta_{\Lambda_{a,b}}(q)$$

**Proof.** We prove this by supposing either

$$\theta_{p,n}(q) = \theta_{p,b}(q) \text{ and } \theta_{p,2m+n}(q) = \theta_{p,2a+b}(q) \quad (7)$$

or

$$\theta_{p,n}(q) = \theta_{p,b+p}(q) \text{ and } \theta_{p,2m+n}(q) = \theta_{p,2a+b+p}(q). \quad (8)$$

From Eq. 7, we have four subcases corresponding to  $n \equiv \pm b \pmod{2p}$  and  $2m + n \equiv \pm(2a + b) \pmod{2p}$ . If  $n \equiv b \pmod{2p}$ , one finds that  $m \equiv a$

mod  $p$  or  $m \equiv -a - b \pmod{p}$ . If  $n \equiv -b \pmod{2p}$ , one finds that  $m \equiv a + b \pmod{p}$  or  $m \equiv -a \pmod{p}$ .

From Eq. 8, we have four subcases as well, corresponding to  $n \equiv \pm(b+p) \pmod{2p}$  and  $2m + n \equiv \pm(2a + b + p) \pmod{2p}$ . If  $n \equiv b + p \pmod{2p}$ , then either  $m \equiv a \pmod{p}$  or  $m \equiv -a - b \pmod{p}$ . And if  $n \equiv -b - p \pmod{2p}$ , then either  $m \equiv a + b \pmod{p}$  or  $m \equiv -a \pmod{p}$ .

Therefore, if  $n \equiv b \pmod{p}$ , then  $m \equiv a \pmod{p}$  or  $m \equiv -a - b \pmod{p}$ . If  $n \equiv -b \pmod{p}$ , then  $m \equiv a + b \pmod{p}$  or  $m \equiv -a \pmod{p}$ .  $\square$

**Remark 2.1.** Notice that in the case of  $p = 2$ , there are 4 lattices  $\Lambda_{a,b}$  corresponding to choices of  $a$  and  $b$  modulo 2. One finds that  $\theta_{\Lambda_{0,1}}(q) = \theta_{\Lambda_{1,1}}(q)$  (which is given as Equation (3.9) in Lemma 3.1 of [2]), so there are 3 associated theta series.

**Remark 2.2.** In the case of  $p = 3$ , among the 9 lattices, one finds that

$$\begin{aligned}\theta_{\Lambda_{0,1}}(q) &= \theta_{\Lambda_{2,1}}(q) = \theta_{\Lambda_{1,2}}(q) = \theta_{\Lambda_{0,2}}(q), \\ \theta_{\Lambda_{1,1}}(q) &= \theta_{\Lambda_{2,2}}(q), \text{ and} \\ \theta_{\Lambda_{1,0}}(q) &= \theta_{\Lambda_{2,0}}(q),\end{aligned}$$

giving a total of 4 associated theta series.

For general odd  $p$ , among the  $p^2$  lattices, there are  $\frac{(p+1)^2}{4}$  associated theta series.

### 3. Theta functions of codes over $\mathcal{R}$

Let  $p \nmid \ell$  and

$$\mathcal{R} := \mathcal{O}_K/p\mathcal{O}_K = \{a + b\omega : a, b \in \mathbb{F}_p, \omega^2 + \omega + d = 0\}.$$

A linear code  $C$  of length  $n$  over  $\mathcal{R}$  is an  $\mathcal{R}$ -submodule of  $\mathcal{R}^n$ . The dual is defined as  $C^\perp = \{u \in \mathcal{R}^n : u \cdot \bar{v} = 0 \text{ for all } v \in C\}$ . If  $C = C^\perp$  then  $C$  is self-dual. We define

$$\Lambda_\ell(C) := \{x \in \mathcal{O}_K^n : \rho_\ell(x) \in C\},$$

where  $\rho_\ell : \mathcal{O}_K \rightarrow \mathcal{O}_K/p\mathcal{O}_K \rightarrow \mathcal{R}$ . In other words,  $\Lambda_\ell(C)$  consists of all vectors in  $\mathcal{O}_K^n$  which when taken mod  $p\mathcal{O}_K$  componentwise are in  $\rho_\ell^{-1}(C)$ . This method of lattice construction is known as Construction A.

For  $0 \leq a, b \leq p - 1$ , let  $r_{a+pb} = a - b\omega$ , so  $\mathcal{R} = \{r_0, \dots, r_{p^2-1}\}$ . For a codeword  $u = (u_1, \dots, u_n) \in \mathcal{R}^n$  and  $r_i \in \mathcal{R}$ , we define the counting function

$$n_i(u) := \#\{i : u_i = r_i\}.$$

The complete weight enumerator of the  $\mathcal{R}$  code  $C$  is the polynomial

$$cwe_C(z_0, z_1, \dots, z_{p^2-1}) = \sum_{u \in C} z_0^{n_0(u)} z_1^{n_1(u)} \dots z_{p^2-1}^{n_{p^2-1}(u)}. \quad (9)$$

We can use this polynomial to find the theta function of the lattice  $\Lambda_\ell(C)$ .

**Lemma 3.1.** *Let  $C$  be a code defined over  $\mathcal{R}$  and  $cwe_C$  its complete weight enumerator as above. Then,*

$$\theta_{\Lambda_\ell(C)}(q) = cwe_C(\theta_{\Lambda_{0,0}}(q), \theta_{\Lambda_{1,0}}(q), \dots, \theta_{\Lambda_{p-1,p-1}}(q))$$

**Proof.** Since

$$\theta_{\Lambda_\ell(C)}(q) = \sum_{z \in \Lambda_\ell(C)} q^{z \cdot \bar{z}},$$

one has

$$\begin{aligned} \theta_{\Lambda_\ell(C)}(q) &= \sum_{u \in C} \theta_{\Lambda_\ell(u)}(q), \\ &= \sum_{u \in C} \sum_{x \in u + p\mathcal{O}_K^n} q^{x \cdot \bar{x}}, \\ &= \sum_{u \in C} \prod_{j=1}^n \sum_{x \in u_j + p\mathcal{O}_K} q^{x \cdot \bar{x}} \text{ (for } u = (u_1, \dots, u_n)), \\ &= \sum_{u \in C} \prod_{j=1}^n \theta_{u_j + p\mathcal{O}_K}(q), \\ &= \sum_{u \in C} \prod_{i=0}^{p^2-1} (\theta_{\tilde{r}_i + p\mathcal{O}_K}(q))^{n_i(u)} \text{ (where } \tilde{r}_{a+pb} = a - b\omega_\ell \in \mathcal{O}_K), \\ &= cwe_C(\theta_{\tilde{r}_0 + p\mathcal{O}_K}(q), \theta_{\tilde{r}_1 + p\mathcal{O}_K}(q), \dots, \theta_{\tilde{r}_{p^2-1} + p\mathcal{O}_K}(q)), \\ &= cwe_C(\theta_{\Lambda_{0,0}}(q), \theta_{\Lambda_{1,0}}(q), \dots, \theta_{\Lambda_{p-1,p-1}}(q)), \end{aligned}$$

which completes the proof.  $\square$

### 3.1. A MacWilliams identity

Let  $C^\perp$  be the dual code to  $C$ . One has the following MacWilliams identity (from Theorem 4.1 of [1]).

**Theorem 3.1.** *Let  $\chi : (\mathcal{R}, +) \rightarrow (\mathbb{C}^*, \times)$  be a character of the additive group of  $\mathcal{R}$  whose restriction to any nonzero left ideal of  $\mathcal{R}$  is nontrivial.*

Then

$$cwe_{\mathcal{C}^\perp}(z_0, \dots, z_{p^2-1}) = \frac{1}{p^2} cwe_{\mathcal{C}}(M(z_0, \dots, z_{p^2-1})),$$

where  $M$  is the matrix defined by

$$M = (\chi(r_i \overline{r_j}))_{0 \leq i \leq p-1, 0 \leq j \leq p-1}.$$

To apply this theorem, we need an appropriate character. Define  $\chi$  by  $\chi(a + b\omega) = e^{2\pi i b/p}$ . Any non-zero ideal  $I \subset \mathcal{R}$  contains an element of  $\mathcal{R} - \{0, 1, \dots, p-1\}$ , so there is some  $a + b\omega \in I$  with  $b \neq 0$ , meaning  $\chi$  acts non-trivially on  $I$ . A calculation shows that

$$(a + b\omega)(\overline{s + t\omega}) = (as - at + btd) + (bt - as)\omega,$$

so  $\chi((a + b\omega)(\overline{s + t\omega})) = e^{(bs - at)2\pi i/p}$ . This is independent of  $d$ , so we obtain the same MacWilliams identity for codes over  $\mathbb{F}_{p^2}$  and  $\mathbb{F}_p \times \mathbb{F}_p$ .

In the case of  $p = 2$ , for example, such identities can be made explicit; see [2] and [1] among others.

### 3.2. A generalization of the symmetric weight enumerator polynomial

In [2], for  $p = 2$ , the symmetric weight enumerator polynomial  $swe_{\mathcal{C}}$  of a code  $\mathcal{C}$  over a ring or field of cardinality 4 is defined to be

$$swe_{\mathcal{C}}(X, Y, Z) = cwe_{\mathcal{C}}(X, Y, Z, Z).$$

For  $\Lambda_{\mathcal{C}}(q)$  the lattice obtained from  $\mathcal{C}$  by Construction A, by Theorem 5.2 of [2], one can then write

$$\theta_{\Lambda_{\mathcal{C}}}(q) = swe_{\mathcal{C}}(\theta_{\Lambda_{0,0}}(q), \theta_{\Lambda_{1,0}}(q), \theta_{\Lambda_{0,1}}(q)).$$

(These theta functions are referred to as  $A_d(q)$ ,  $C_d(q)$ , and  $G_d(q)$  in [2] and [10].)

For  $p > 2$ , however, there are  $\frac{(p+1)^2}{4}$  (which is larger than 3) theta functions associated to the various lattices, so our analog of the symmetric weight enumerator polynomial needs more than 3 variables.

**Example 3.1.** For  $p = 3$ , from Remark 2.2, we have four theta functions corresponding to the lattices  $\Lambda_{a,b}$ , namely

$$\theta_{\Lambda_{0,0}}(q), \theta_{\Lambda_{1,0}}(q), \theta_{\Lambda_{1,1}}(q), \theta_{\Lambda_{0,1}}(q).$$

If we define the ‘‘symmetric weight enumerator for  $p = 3$ ’’ to be

$$swe_{\mathcal{C}}(X, Y, Z, W) = cwe_{\mathcal{C}}(X, Y, Y, Z, W, Z, Z, Z, W),$$

then one finds that

$$\theta_{\Lambda_\ell C}(q) = cwe_C(\theta_{\Lambda_{0,0}}(q), \theta_{\Lambda_{1,0}}(q), \dots, \theta_{\Lambda_{2,2}}(q)), \quad (10)$$

$$= swe_C(\theta_{\Lambda_{0,0}}(q), \theta_{\Lambda_{1,0}}(q), \theta_{\Lambda_{1,1}}(q), \theta_{\Lambda_{0,1}}(q)). \quad (11)$$

**Problem 1.** Define an symmetric weight enumerator, analogous to the  $p = 2$  case, for codes defined over  $\mathcal{R}$  for  $p > 3$ . Write a MacWilliams identity for the symmetric weight enumerator and determine an explicit relation between the symmetric weight enumerator and theta functions.

#### 4. The injectivity of Construction A.

For a fixed prime  $p$ , let  $\mathcal{R} = \mathcal{O}_K/p\mathcal{O}_K$  and  $C$  be a linear code over  $\mathcal{R}$  of length  $n$  and dimension  $k$ . An admissible level  $\ell$  is an integer  $\ell$  such that  $\mathcal{R}$  is isomorphic to  $\mathbb{F}_{p^2}$  or  $\mathbb{F}_p \times \mathbb{F}_p$ . For an admissible  $\ell$ , let  $\Lambda_\ell(C)$  be the corresponding lattice as in the previous section. Then, the **level  $\ell$  theta function**  $\theta_{\Lambda_\ell(C)}(\tau)$  of the lattice  $\Lambda_\ell(C)$  is determined by the complete weight enumerator  $cwe_C$  of  $C$ , evaluated on the theta functions defined on cosets of  $\mathcal{O}_K/p\mathcal{O}_K$ . We consider the following questions:

- i) How do the theta functions  $\theta_{\Lambda_\ell(C)}(\tau)$  of the same code  $C$  differ for different levels  $\ell$ ?
- ii) Can non-equivalent codes give the same theta functions for all levels  $\ell$ ?

Next we see how this can be made explicit for the case  $p = 2$ .

##### 4.1. The characteristic $p = 2$

For  $p = 2$  case these questions are fully answered in [10]. We have the following:

**Theorem 4.1** ([10], **Thm. 1**). Let  $p = 2$  and  $C$  be a code defined over  $\mathcal{R}$ . For all admissible  $\ell, \ell'$  such that  $\ell > \ell'$ , the following holds

$$\theta_{\Lambda_\ell}(C) = \theta_{\Lambda_{\ell'}}(C) + \mathcal{O}(q^{\frac{\ell'+1}{4}}).$$

**Theorem 4.2** ([10], **Thm. 2**). Let  $p = 2$  and  $C$  be a code of length  $n$  defined over  $\mathcal{R}$  and  $\theta_{\Lambda_\ell}(C)$  be its corresponding theta function for level  $\ell$ . Then the following hold:

- i) For  $\ell < \frac{2(n+1)(n+2)}{n} - 1$  there is a  $\delta$ -dimensional family of symmetrized weight enumerator polynomials corresponding to  $\theta_{\Lambda_\ell}(C)$ , where  $\delta \geq \frac{(n+1)(n+2)}{2} - \frac{n(\ell+1)}{4} - 1$ .

- ii) For  $\ell \geq \frac{2(n+1)(n+2)}{n} - 1$  and  $n < \frac{\ell+1}{4}$  there is a unique symmetrized weight enumerator polynomial which corresponds to  $\theta_{\Lambda_\ell}(C)$ .

These results were obtained by using the explicit expression of theta in terms of the symmetric weight enumerator valuated on the theta functions of the cosets. Hence, a solution to Problem 1 most likely would lead to obtaining such results for all  $p > 2$  and admissible  $\ell$ . In this paper we use the complete weight enumerator polynomial to get similar results.

#### 4.2. The characteristic $p > 2$

Let  $C$  be a code defined over  $\mathcal{R}$  for a fixed  $p > 2$ . Let the complete weight enumerator of  $C$  be the degree  $n$  polynomial

$$cwe_C = f(x_0, \dots, x_r)$$

for  $r = p^2 - 1$ . Then from Lemma 3.1 we have that

$$\theta_{\Lambda_\ell(C)}(\tau) = f(\theta_{\Lambda_{0,0}}(\tau), \dots, \theta_{\Lambda_{p-1,p-1}}(\tau))$$

for a given  $\ell$ . First we want to address how  $\theta_{\Lambda_\ell(C)}(\tau)$  and  $\theta_{\Lambda_{\ell'}(C)}(\tau)$  differ for different  $\ell$  and  $\ell'$ . We have the following:

**Theorem 4.3.** *Let  $C$  be a code defined over  $\mathcal{R}$ . For all admissible  $\ell, \ell'$  the following holds*

$$\theta_{\Lambda_\ell(C)} - \theta_{\Lambda_{\ell'}(C)} = \sum_{i=0}^s a_i q^s$$

for some  $a_i \in \mathbb{Z}$  and  $s \in \mathbb{Z}^+$ .

**Corollary 4.1.** *Let  $p$  be a fixed prime and  $\ell$  any square free integer such that  $K = \mathbb{Q}(\sqrt{-\ell})$  and  $\mathcal{R} := \mathcal{O}_K/p\mathcal{O}_K$  is isomorphic to  $\mathbb{F}_{p^2}$  or  $\mathbb{F}_p \times \mathbb{F}_p$ . For a given code  $C$  defined over  $\mathcal{R}$ , the theta series  $\theta_{\Lambda_\ell(C)}$  is the same for almost all  $\ell$ .*

**Theorem 4.4.** *Let  $C$  be a code of size  $n$  defined over  $\mathcal{R}$  and  $\theta_{\Lambda_\ell(C)}$  be its corresponding theta function for level  $\ell$ . Then, for large enough  $\ell$ , there is a unique complete weight enumerator polynomial which corresponds to  $\theta_{\Lambda_\ell(C)}$ .*

The proofs of Theorems 4.3 and 4.4 are provided in [11] where explicit bounds for  $\ell$  are provided for small  $p$ .

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